

Effects of Relay Selection Strategies on the Spectral Efficiency of Wireless Systems with Half- and Full-duplex Nodes

Carlos H. M. de Lima *Member, IEEE*, Hirley Alves *Member, IEEE*, Pedro. H. J. Nardelli, and Matti Latva-aho *Senior Member, IEEE*

Abstract

This work proposes an analytical framework to study how relay selection strategies perform in half- and full-duplex deployments by combining renewal theory and stochastic geometry. Specifically, we assume that the network nodes – operating in either half- or full-duplex mode – are scattered according to a two-dimensional homogeneous Poisson point process to compute the relay selection cost by using a semi-Markov process. Our results show: *(i)* fixed relay outperforms the reactive option in either cases, *(ii)* the performance of both reactive and fixed relay strategies depends on the self-interference attenuation in full-duplex scenarios, evincing when they outperform the half-duplex option, and *(iii)* the reactive relay selection suffers from selecting relays at hop basis, while the fixed relay selection benefits most from the full-duplex communication.

I. INTRODUCTION

In the recent years, Full-Duplex (FD) communication has gained considerable attention from both academy and industry [1], [2]. By allowing simultaneous transmission and reception on the same frequency band, FD networks can potentially double the spectral efficiency compared to

C. H. M. de Lima is with São Paulo State University (UNESP), São João da Boa Vista, Brazil. Email: carlos.lima@sjbv.unesp.br.

H. Alves, P. H. J. Nardelli and M. Latva-aho are with the Centre for Wireless Communications (CWC), University of Oulu, Finland. E-mail:firstname.lastname@oulu.fi.

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current Half-Duplex (HD) schemes. However, such potential is harmed by Self-Interference (SI) from the transmit to receive antenna [1], [2]. Due to recent advances in antenna design associated with analog and digital interference cancellation, most of the SI can be mitigated, and therefore FD communication becomes feasible and are a step towards meeting high demands of spectral efficiency of 5G [3].

Current works have shown that FD is a viable solution for small cell deployments [2], due to not only cost and size constraints but also to compatibility with current HD systems. Then, due to the FD capability, the base station can, for instance, simultaneously schedule uplink and downlink transmissions, see for instance [4]. Moreover, base station can also act as relays for the legacy network as well as for its own subscribed users. Cooperative diversity appeared as a way to combat the effects of the fading, by allowing a single antenna user to experience spatial diversity [5]. Such schemes perform even better when the relay operates in a FD fashion [6].

Bearing this in mind, and in the context of a dense deployment of small cells, relaying selection become an attractive solution given a trade-off between complexity and efficiency [2], [7]. Relay selection was initially proposed in [7] for HD cooperative networks and has gained considerable attention since. Nonetheless, most of the works so far have focused on HD selection algorithms, except for [8], [9] where the authors focus on the performance of a three-node relaying scenario. For instance, [8] focus on selection algorithms for FD amplify-and-forward and model the residual SI at relay as a Rayleigh random variable. While in [9], the authors address the performance of a FD decode-and-forward protocol and assume simpler residual SI model (constant attenuation factor).

In this work, we assess the performance at the network level of a relay selection procedure for cooperative diversity protocols in FD mode. To do so, we model the relaying selection procedure as a semi-Markov process and then we investigate the impact of the SI and Co-Channel Interference (CCI) on the network performance. Moreover, a dynamic relay selection procedure is considered where a suitable relay is selected at each hop and the cost of this selection procedure is incorporated into the achievable rate. And yet, we resort to stochastic geometry to model network deployments and capture the dynamics of the network, more specifically the aggregate interference at the receiver, which can be characterized through a cumulant-based framework introduced in [10], [11].

Our main contributions are summarized as follows:

- The standard binary tree algorithm is used to implement a relay selection procedure for FD networks;
- The framework originally introduced in [12] is extended to account for the FD relaying operation;
- The performance of the aforesaid relay selection protocol are evaluated and compared to its HD counterpart;
- Different from [8], [9], our network level analysis accounts for a more detailed model of the residual SI based on [13] and also characterizes the CCI at the receiver. Actually, instead of modeling the residual SI by a constant value, we consider a more detailed model through a random variate drew from a Ricean distribution with large K -factor.

The remainder of this paper is organized as follows: Section II describes the problem under investigation and introduces the system model. Section III describes the network topology and introduces our analytical framework based on the stochastic geometry, while Section IV characterizes the relay selection procedures and the probability generating function concept. Section V introduces the semi-Markov process that is used to evaluate the performance of the proposed relay selection mechanism. In Section VI, performance metrics are introduced and then utilized to obtain the numerical results and carry out their discussion. We draw conclusions and final remarks in Section VII.

II. PROBLEM DESCRIPTION AND SYSTEM MODEL

As aforesaid, full duplex networks constitute a promising alternative to achieve high data rates in the upcoming 5G systems. In such networks, cooperative nodes benefit most from the inherent rate increase allowed by transmitting and receiving simultaneously. Typically, available literature concentrates on assessing the throughput that is attainable by combining full duplex and cooperative diversity protocols, and overlooks the impact of the underlying relay selection on the final achievable data rate [8].

In this paper, we contribute to previous results by introducing an analytical framework that combines stochastic geometry and semi-Markov processes to account for the cost of selecting a suitable relay. On that regard, we initially use stochastic geometry to derive the distributions of the received power and aggregate interference experienced by the receiver of interest in Section III. Thereafter, we use the Probability Generating Function (PGF) concept [14], [15] to characterize

the Contention Resolution Interval (CRI) (holding time) of the relay selection algorithms as can be seen in Sections IV and V. Our investigations extend [12] in which Authors introduce an analytic framework based on semi-Markov processes with constant holding time and apply it to study an one-dimensional network model.

We assess the performance of interference networks operating in FD mode under composite fading channel – Lognormal shadowing and Nakagami- m fading. The Downlink (DL) of a traditional HD network constitutes our benchmark scenario wherein the user of interest is interfered by surrounding small cells. Base stations independently schedule a random user terminal in every transmission interval. All communicating nodes use antennas with omnidirectional radiation pattern. Base stations and User Equipments (UEs) are also assumed to have full buffer and symmetric traffic patterns [16].

Consider all channels are quasi-static and undergo a composite fading distribution [17], which is approximated by a lognormal distribution with mean and variance (in logarithmic scale) given by $\mu_{\text{dB}} = \xi [\psi(m) - \ln(m)] + \mu_{\Omega_p}$ and $\sigma_{\text{dB}}^2 = \xi^2 \zeta(2, m) + \sigma_{\Omega_p}^2$, where $\psi(m)$ is the Euler psi function and $\zeta(2, m)$ is the generalized Riemann zeta function [18]. Radio links are affected by path-loss, large-scale shadowing and multi-path fading which are assumed to be mutually independent and multiplicative phenomena [19]. The received power at the user of interest u_0 from an arbitrary transmitter b_i located r_{i0} meters away is $Y_{i0} = p_{i0} r_{i0}^{-\alpha} x_{i0}$, where p_{i0} yields the transmit power, α is the path-loss exponent and x_{i0} represents the composite fading channel. In what follows, when there is no fear of ambiguity let us omit the subscripts of the interference component terms. Fig. 1 illustrates a realization of the random network topology where UEs and small cell Base Stations (BSs) are uniformly scattered over network area.

III. NETWORK DEPLOYMENT AND STOCHASTIC GEOMETRY

Active nodes constitute a homogeneous Poisson Point Process (PPP) Φ with intensity λ in \mathbb{R}^2 . The number of nodes in an arbitrary region \mathcal{R} of area A is a Poisson Random Variable (RV) with parameter λA [20].

The fading effect is assumed as a random mark associated with each point of Φ , thus $\tilde{\Phi} = \{(\varphi, x); \varphi \in \Phi\}$ corresponds to a marked point process on the product space $\mathbb{R}^2 \times \mathbb{R}^+$, whose random points φ denoting transmitters locations and belong to Φ [20], [21]. The network deployment model is given by a spatial PPPs Φ^{BS} (Φ^{UE}), whose random points φ represent

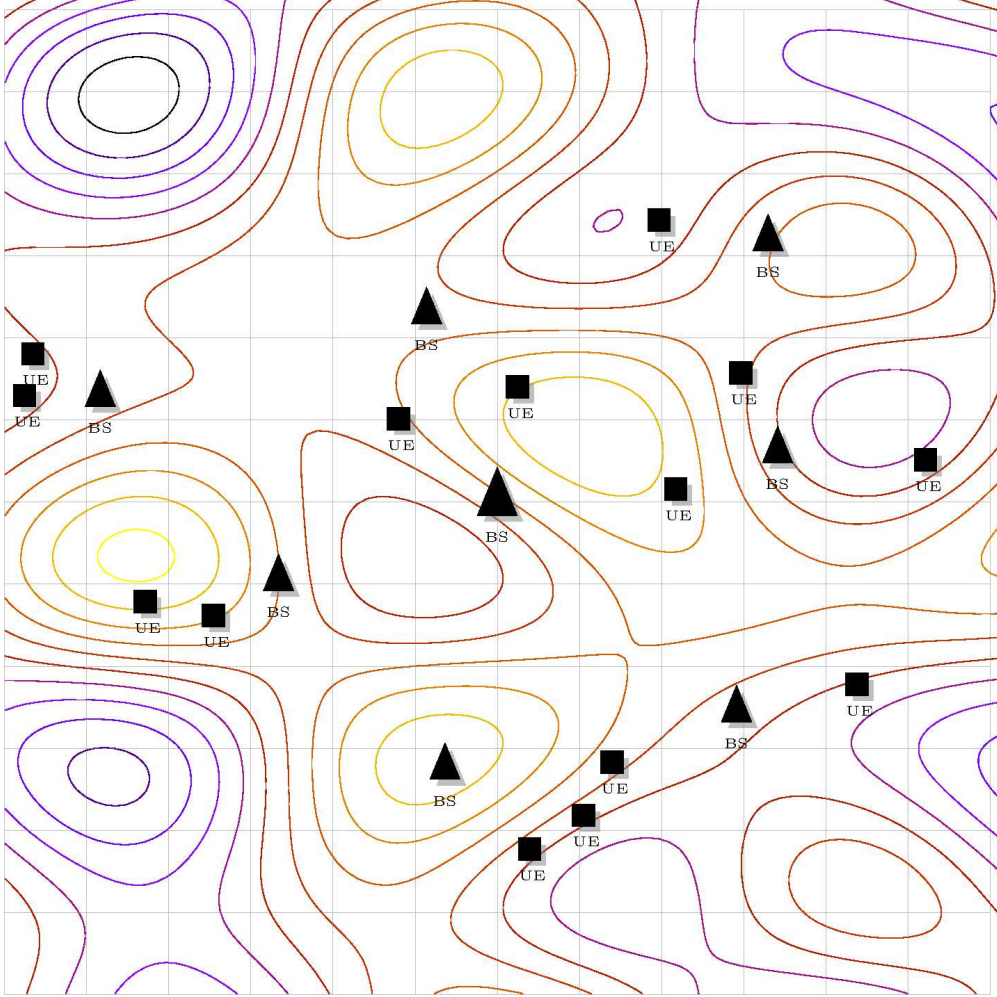


Fig. 1. Illustration of a random (bi-dimensional) deployment of base stations and user terminals over an arbitrary network area. Shaded squares represent user terminals, while shaded up-triangles depict base stations. A heat map represents the corresponding random composite fading channel where the fading intensity varies from red/strong to dark blue/weak.

the locations of BSs (UEs). In addition, we consider a fixed number of potential relays that detect source's message and take part in the selection procedure, it is well known that such conditional distribution of points within an arbitrary forwarding region A is described by a binomial point process [21]. As shown next, by employing stochastic geometry concepts to this deployment model, it is possible to conveniently derive closed form expressions to characterize the distributions of the received power and corresponding aggregate CCI with respect to the receiver of interest. Indeed, the path loss attenuation (function of the distance to the tagged user), the lognormal shadowing and Nakagami- m fading are combined using the framework of

[10], [22] to obtain the distribution of the received power at the user of interest.

A. Received Power from a Random Transmitter

We resort to the framework introduced in [10], [11] to attain the characteristic function and the respective cumulants (κ) of the received power and aggregate interference of random transmitters within the forwarding region of the user of interest. Consider a RV $Y = R^{-\alpha}X$ describing the power received at the tagged receiver from a random transmitter in \mathcal{O} (as defined in Section II), with R varying from R_m to R_M , and $X \sim \text{Lognormal}(\mu_{\text{dB}}, \sigma_{\text{dB}})$ representing the squared-envelope of the composite fading channel. Then, the characteristic function of Y is [11, Appendix B] $\Psi_Y(\omega) = 2(R_M^2 - R_m^2)^{-1}E_X[\mathcal{R}(\omega)]$ where $\mathcal{R}(\omega) = \int_{R_m}^{R_M} \exp(j\omega pr^{-\alpha}x)rdr$ and $E_X[\cdot]$ yields the expectation of the enclosed expression over the RV X , and is a general formulation which characterizes the distribution of any random transmitter within the reception range of a tagged receiver. Thereafter, the n th cumulant of Y is given by [11, Appendix C]

$$\kappa_n = \frac{1}{\imath^n} \sum_{k=0}^n g^{(k)}(\beta_0) \cdot B_{n,k}[\beta_1, \beta_2, \dots, \beta_{(n-k+1)}], \quad (1)$$

where \imath yields the imaginary unity, $g(u) = \ln(u)$, $B_{n,k}[\beta_1, \beta_2, \dots, \beta_{(n-k+1)}]$ is the partial Bell polynomial [23] and $\beta_n = \imath^n p^n \times \frac{R_m^{2-n\alpha} - R_M^{2-n\alpha}}{n\alpha - 2} E_X[x^n]$.

Next, we use the aforesaid cumulant-based framework to characterize the aggregate interference at the user of interest for both the HD and FD configuration scenarios. Under the assumptions of Section II and with respect to the tagged receiver, we now use our analytical framework to characterize the resulting aggregate CCI in two Evaluation Scenarios (ESs): ES1 is the benchmark scenario where BSs and UEs operate in HD mode; while in ES2 both BSs and UEs operate in FD mode.

B. Aggregate CCI from a Poisson Field of Interferers

In our benchmark scenario, denoted *HD relaying*, all nodes operate in HD mode. The corresponding aggregate CCI at a tagged receiver is then given by $Z_0^{\text{HD}} = \sum_{(\varphi_i, x_i) \in \tilde{\Phi}^{\text{BS}}} Y_{i0}$. Notice that the tagged receiver only experiences interference coming from nearby base stations since we assume the DL as the benchmark scenario.

In this context the aggregate CCI at the tagged receiver is [10, Section V]

$$\kappa_n(\tilde{\Phi}^{\text{BS}}) = \frac{2\pi\lambda p^n}{n\alpha - 2} (R_m^{2-\alpha n} - R_M^{2-\alpha n}) E_X^n[0, \infty]. \quad (2)$$

C. Aggregate CCI in the FD Configuration

In the FD configuration, the tagged receiver is subject to the interference from FD serving base stations, other FD UEs and its intrinsic SI component. The self-interference channel between transmitting and receiving antennas of a full duplex transceiver exhibits a strong line of sight component and can be represented using a Ricean distribution with large K -factor [24] (a K -factor of 14.8dB is used, which corresponds to a Nakagami distribution with the parameter $m = 16$). As a result, the aggregate CCI at the tagged receiver is characterized by

$$Z_0^{\text{FD}} = \delta p_{00} x_{00} + \sum_{(\varphi_i, x_i) \in \tilde{\Phi}^{\text{BS}}} Y_{i0} + \sum_{(\varphi_j, x_j) \in \tilde{\Phi}^{\text{UE}}} Y_{j0}, \quad (3)$$

where p_{00} and δ represent the SI component and the corresponding attenuation factor, respectively.

Since BSs and UEs are assumed to be independently scattered over the network deployment area, the resulting processes (Φ^{BS} and Φ^{UE} , respectively) from each such tier are also independent [20]. Therefore, the cumulants additivity property is employed to obtain the aggregate CCI from multiple interfering tiers [10]. In other words, if U and V are independent RVs, then, we can derive the cumulant of their sum as $\kappa_n(U + V) = \kappa_n(U) + \kappa_n(V)$.

D. Signal-to-Interference Ratio and Outage Probability

By using the results from previous sections, it is straightforward to compute the Signal-to-Interference Ratio (SIR) and outage probability of the evaluation scenarios under study. The SIR and outage probability are used to assess how the relay selection strategies perform in HD and FD configuration mode as follows.

Proposition 1: Let V_0 and V be Normal RVs (in logarithmic scale) representing the power received from the desired transmitter and the CCI at the receiver of interest. Under the assumption of the composite fading with Gamma-lognormal distribution, the SIR at the tagged receiver is $\Gamma \sim \text{Normal}(\mu_{V_0} - \mu_V, \sigma_{V_0}^2 + \sigma_V^2)$, and the outage probability is given by $\Pr[\Gamma < \gamma_{\text{th}}] = \text{Q}[(\mu_{\Gamma} - \gamma_{\text{th}})/\sigma_{\Gamma}]$, where $\mu_{\Gamma} = \mu_{V_0} - \mu_V$ and $\sigma_{\Gamma}^2 = \sigma_{V_0}^2 + \sigma_V^2$. Notice that the lognormal parameters

can be attained as follows $\mu = \ln\left(\frac{\kappa_1^2}{\sqrt{\kappa_1^2 + \kappa_2}}\right)$, and $\sigma^2 = \ln\left(1 + \kappa_2/\kappa_1^2\right)$, where κ_1 and κ_2 are given by (1).

Proof: The SIR distribution is given by the quotient of two independent lognormal RVs, namely, e^{V_0} which is the received power from the target transmitter, and e^V which is an equivalent lognormal RV approximating the aggregate CCI at the tagged receiver [25]. ■

IV. RELAY SELECTION ALGORITHM AND PROBABILITY GENERATING FUNCTION

Hereafter, we characterize the cost of selecting a suitable relay at hop basis on the achievable throughput. To do that, we first account for the relay selection overhead using the PGF of the CRI and later combine the impact of the network characteristics with the relay selection mechanism throughout a semi-Markov process. In this work, a totally random approach based solely on the standard tree algorithm is used to implement the relay selection algorithm [26]. It is worth noticing we use the conditional PGF of the CRI length (L_N) that is derived in [14], [15] for the same standard tree algorithm under investigation. Next, due to space limitations, we only summarize the main results regarding the relay selection procedure.

A. Relay Selection Algorithm

Two relay selection approaches are considered in this work: (i) a static one whereby a relay node is preassigned to forward source packets; and (ii) a reactive strategy by which a relay node is identified based on the selection criteria – *i.e.*, the neighboring node that provides the longest advancement towards the final destination at each hop. While the former configuration may be more applicable to legacy networks where a smallcell may provide the last hop to cell edge users, the latter is tailored to upcoming 5G networks where humans and machines communicate in large scale deployment scenarios [27]. On one hand, the static relay procedure relies on long term characteristics of the deployment scenarios, *i.e.* the relay is selected a priori following a network planning (it is arbitrary in our investigations) and does not change. On the other hand, the reactive relay procedure selects a random (new) relay within the transmission range every time a packet needs to be forwarded to the destination [12], [14], [28].

Succinctly, the reactive relay selection algorithm is implemented through a carrier sense collision avoidance mechanism similar to the IEEE 802.11 handshake using Request to Send (RTS)/Clear To Send (CTS) messages. The source transmits a RTS packet to initiate the relay

selection procedure. Nodes that listen to this request reply with a CTS packet based on the predetermined probability of accessing the channel, p . If a collision occurs, nodes that have transmitted in previous slot retransmit or not based on a random access process similar to a Q -sided coin. To select the best suitable relay (greed criteria), *i.e.* the closest node to the destination whether there is one available, the source node should receive the replies from all the candidate relays [14], [15]. We consider the conditional CRI length when N nodes initially collide (which corresponds to a binomial point process [21]). Here, aiming at shorten the selection procedure¹, the first node (at a random location within transmission range) to successfully reply to the source is selected as the next hop relay – note that this mechanism differs from the one used in [14], [15] which waits for the replies of all neighboring relay candidates in order to select the most suitable based on the predefined selection criteria.

We use the relay selection mechanism introduced in [14], [15] which combines the standard tree algorithm to resolve contentions and sectoral decision regions to identify the potential relays. Fig. 2 illustrates a snapshot of the sectoral decision region that is used in conjunction with the standard tree-based solution. In the next section, the PGF is used to characterize the distribution of the CRI associated with the selection of a random relay at hop basis.

Remark 1: Notice that current 802.11 standards implements the RTS/CTS handshake as an optional feature to avoid frames collisions; besides, more recent standards such as 802.22 and 802.11af also incorporate cognitive radio functionalities which makes it even easier to deal with side-information [14].

Remark 2: In LTE systems, the contention-based random channel is based on ALOHA-type communication and is susceptible to performance degradation due to a high collision probability in the transmission of the preambles in dense deployments. Indeed, the contention-based operation of the LTE Physical Random Access Channel is also established by the signaling exchange of messages, namely, preamble transmission, random access response, connection request and contention resolution [27], [29].

¹At the cost of selecting a relay that is not necessarily the closest one to the final destination

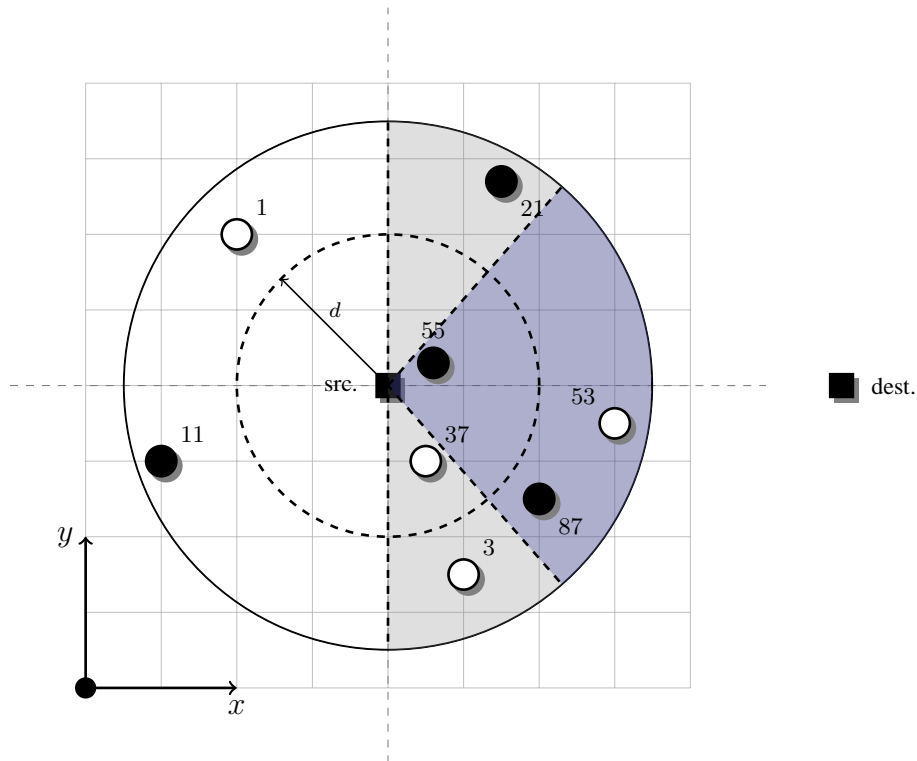


Fig. 2. Illustration of the sectoral decision region with $Q = 2$ splitting groups and angular aperture of 180 degrees. Dashed lines defined the forwarding region (shaded regions in light gray and blue), black circles identify awake nodes and white circles identify asleep nodes. All awake neighbors within the shaded region are eligible relays.

B. PGF of the Relay Selection Algorithm

The conditional CRI length considering a Q -sided fair coin is generalized as,

$$L_N = \begin{cases} 1 & , \text{ if } N = 0, 1; \\ 1 + \sum_{j=1}^Q L_{I_j} & , \text{ if } N \geq 2, \end{cases} \quad (4)$$

where I_j is a discrete RV describing the number of candidate relays that tossed the j value of the Q -side coin. For a binary splitting tree, the PGF of the time spent by a random node in the relay selection procedure is give by the following proposition.

Proposition 2: Under the assumptions of a binary standard tree algorithm (unbiased coin) and

N conflicting nodes, the PGF of the CRI for a tagged packet is

$$G_{N+1}(z) = \frac{1}{2} \left[G_{N+1}^{(0)}(z) + G_{N+1}^{(1)}(z) \right] \quad \text{with} \quad (5)$$

$$G_{N+1}^{(0)}(z) = z \sum_{k=0}^N B_{N,k} G_{k+1}(z) \quad \text{and} \quad (6)$$

$$G_{N+1}^{(1)}(z) = z \sum_{k=0}^N B_{N,k} Q_k(z) G_{N-k+1}(z), \quad (7)$$

where $G_1 = z$, $G_{N+1}^{(0)}(z)$ accounts for the CRI when the tagged node joins the first subset (flipped 0), $G_{N+1}^{(1)}(z)$ describes the resolution interval when the tagged node joins the second subset (flipped 1), and $B_{N,k} = \Pr[k \text{ flipped } 0 \mid N \text{ contending nodes}]$. Note that the tagged packet may be in each subset with equal probability.

Proof: Let $N > 1$ be the number of contending nodes that start the CRI. For a fair binary splinting tree we have $B_{N,k} = \binom{N}{k} 2^{-k}$. Thus, if the first splitting group has k members, the second contains the remaining $N - k$ nodes. Each such sub-CRI is statically indistinguishable from a general CRI with equal number of elements. Therefore, the recursive formula $Q_N(z) = \sum_{k=0}^N B_{N,k} Q_k(z) Q_{N-k}(z)$ is obtained. By the same line of reasoning, we can defined

$$G_{N+1}(z) = \Pr[t_0 = k \mid N \text{ other nodes collide}] z^k; \quad (8)$$

$$G_{N+1}^{(s)}(z) = \Pr[t_0 = k \mid N \text{ other nodes collide} \\ \text{and the tagged node flips } s] z^k, \quad (9)$$

where t_0 yields the time a tagged packet spends in the CRI before being successfully transmitted. Finally, computing the corresponding probabilities for a fair binary spiting tree, we obtain (5). ■

In other words, Proposition 2 gives the time a tagged packet takes to be successfully transmitted. When a random node within source's transmission range is selected, its distance distribution is characterized by (1). To determine the holding time matrix, we resort to the factorial moment concept so as to compute the mean and variance of CRI for a tagged packet from Proposition (2).

Definition 1: Let X be a discrete RV taking non-negative integer values; then, the k th factorial moment of X is given as $E[X(X-1)\cdots(X-k+1)] = G^{(k)}(1^-)$, $k \geq 0$ [30].

Then, we obtain the following results from Definition 1 as $G(1) = \sum_k \Pr[X = k]$, $G'(1) = E[X]$ and $G''(1) = E[X(X - 1)]$.

V. SEMI-MARKOV PROCESS

Markov processes permit represent random systems whose outcome at any given instant depends only on the outcome that proceeds it (memoryless property, if the present is specified, the past has no influence in the future). A Markov chain is special kind of Markov process wherein the future evolution of the process depends on the present state and not on how it arrived at that state [31]. Typically, such Markov models have the property that a transition occurs at every instant – even if the transition returns to the previous occupied state, it occurs anyway. In this work, we use a semi-Markov process that is associated with the embedded Markov chain in Fig. 3 that tracks the relay selection protocols states. The semi-Markov process theory allow us to obtain the steady-state average delay and throughput during a cycle taking any state of the chain as a reference [30], [31]. Hence, at transition instants semi-Markov process behaves just as a Markov process. is used to combine the network topology (including radio channel) that are captured by the stochastic geometry with the routing protocol dynamics that are captured by the PGF framework. With semi-Markov process, the elapsed time between transitions may take several unit intervals, while at transition instants it behaves as a typical Markov process [30]. The semi-Markov process depends on the transition that occurred, though the transition instants follow a distinct probabilistic mechanism [30]. Throughout our investigation we consider the states presented in Fig. 3: s_1 , s_2 and s_3 corresponds to the transmission, relay and re-transmission states, respectively. The corresponding state transition matrix is given by

$$\mathbf{P} = \begin{bmatrix} p_{SD} & (1 - p_{SD})p_{SR} & (1 - p_{SD})(1 - p_{SR}) \\ 1 & 0 & 0 \\ p_{SD} & (1 - p_{SD})p_{SR} & (1 - p_{SD})(1 - p_{SR}) \end{bmatrix} \quad (10)$$

where each element p_{ij} yields the transition probability from state i to j , as well as p_{SD} and p_{SR} yield the probability that the source message is correctly received by the final destination and relay node, respectively.

Remark 3: The Markov process model has the property that a transition is made at every time instant. Differently, the semi Markov model is a more general class of processes where the elapsed time between transitions may take several unit time intervals. Then, it brings much

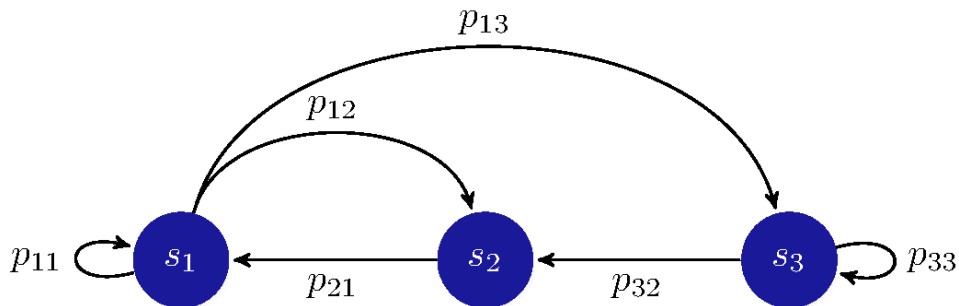


Fig. 3. Illustration of a Semi-Markov process with three states s_i , $i \in \{1, 2, 3\}$. s_1 represents transmission, s_2 relay and s_3 retransmission.

more flexibility to the problem of modeling dynamic probabilistic systems. In fact, the successive state occupancies are governed by the transitions probabilities of an embedded Markov process, while the holding time in each state follows an integer-valued random distribution – the conflict resolution interval, see (4) in Section IV – which depends on both the current and the next states [30].

The PGF of the CRI $G_N(z)$ characterizes the relay selection overhead at each hop. The probability that a certain CRI takes k slots can be retrieved from the respective PGF using the following proposition.

Proposition 3: The probability mass function of L_N , the length of the contention resolution interval (in transmission slots) conditional on the number of contending nodes is $\Pr[X = x] = G^{(x)}(0)/x!$.

Proof: From the PGF definition it is straightforward that

$$G_N(z) = \mathbb{E}[z^X] = \sum_{x=0}^{\infty} \Pr[X = x]z^x. \quad (11)$$

Then, taking the n th derivative of (11) and making $z = 0$ we attain the result in Proposition 3. ■

The holding time matrix incorporates into the embedded Markov chain of Fig. 3 the time spent in a given state between successive transitions.

$$\mathbf{H} = \begin{bmatrix} 1 & 1 + L_N & 1 + L_N \\ 1 & 1 & 1 \\ 1 & 1 + L_N & 1 + L_N \end{bmatrix}, \quad (12)$$

where each element h_{ij} yields the time elapsed in state i before transition to j .

A. Steady – State Throughput Efficiency

Let i be an arbitrary state chosen as reference, and π_i be its steady-state probability where the vector π corresponds to the steady-state distribution of the embedded Markov chain (see Fig. 3). The time delay associated to the transition from state i to j is given by D_{ij} . Similarly, R_{ij} is the reward associated with the transition from state i to j . A cycle is defined as the time between two consecutive passages for the reference state states. A message is successfully delivered to the destination node when the process returns to the state s_1 . The reward function $R(\tau)$ corresponds to the cumulative reward after a period τ and yields the total amount of correct receptions. According to the fundamental renewal-reward theorem [30], the long-term throughput is given by the following value, It is worth noticing that the resulting steady-state throughput efficiency is independent of the choice of the reference state. Moreover, the long term evaluation of $R(\tau)/\tau$ actually corresponds to the normalized throughput of the protocol as introduced in [32] and revisited in [28]

$$\eta = \lim_{\tau \rightarrow \infty} \frac{R(\tau)}{\tau} = \frac{\sum_{i=1}^K \pi_i \bar{R}_i}{\sum_{i=1}^K \pi_i \bar{D}_i}, \quad (13)$$

where K yields the number of states, $\bar{R} = \sum_j^K p_{ij} R_{ij}$ is the mean rewarding, while $\bar{H} = \sum_j^K p_{ij} h_{ij}$ is the mean waiting time.

VI. NUMERICAL RESULTS

For both desired link and interferers, a composite fading channel model with Nakagami- m parameter of $m = 16$ (which corresponds to a Ricean fading channel with factor $K = 14.8$ dB) and shadowing standard deviation of 10 dB are considered as well as channel reciprocity between a communication pairs. Interferes are scattered over the network area with $R_m = 25$ and $R_M = 500$ meters, while $\lambda = 5 \times 10^{-5}$ nodes/m² is considered (which corresponds to nearly 40 interferers in average) and the path loss attenuation is set to 3. The CRI is conditioned on the mean number of candidate relays on the forward region of the source node. Without loss of generality, we consider three contending relays within source's forwarding region to compute the relay selection interval. We assume that data and control information occupies one transmission interval.

In a HD configuration, Fig. 4 shows the steady-state throughput efficiency $\bar{\eta}$ for increasing separation distance between source–destination pair. As evidenced by the reactive relay curve

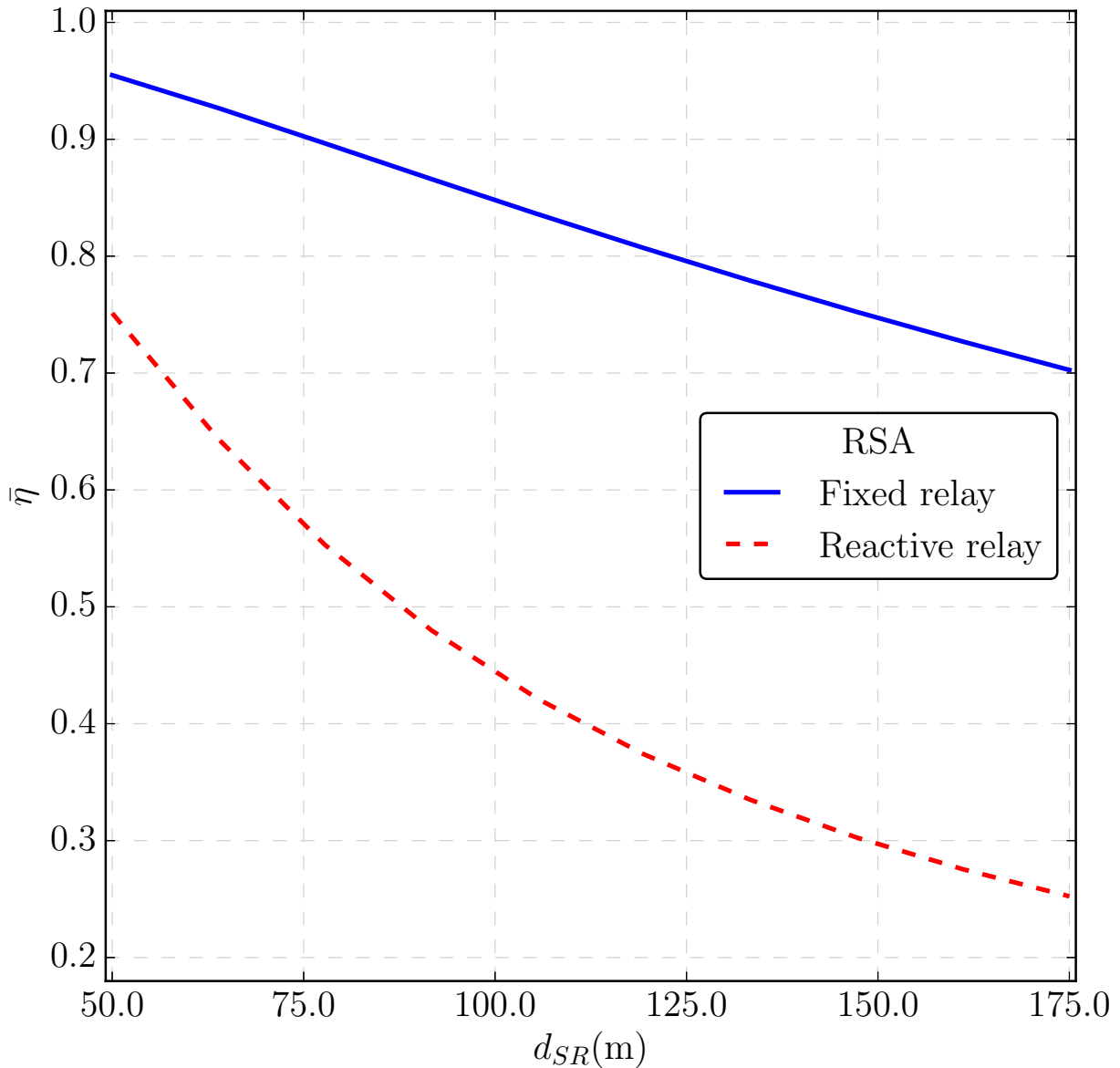


Fig. 4. Steady-state throughput efficiency for increasing separation distance between source – destination pair in HD configuration.

(red dashed line), the steady-state throughput is severely compromised by the relay selection. On the other hand, when a fixed relay is considered within source's transmission range, $\bar{\eta}$ becomes much less susceptible to the degrading effect of longer separation distance between source and destination.

Fig. 5 presents the steady-state throughput for increasing transmission power of source node,

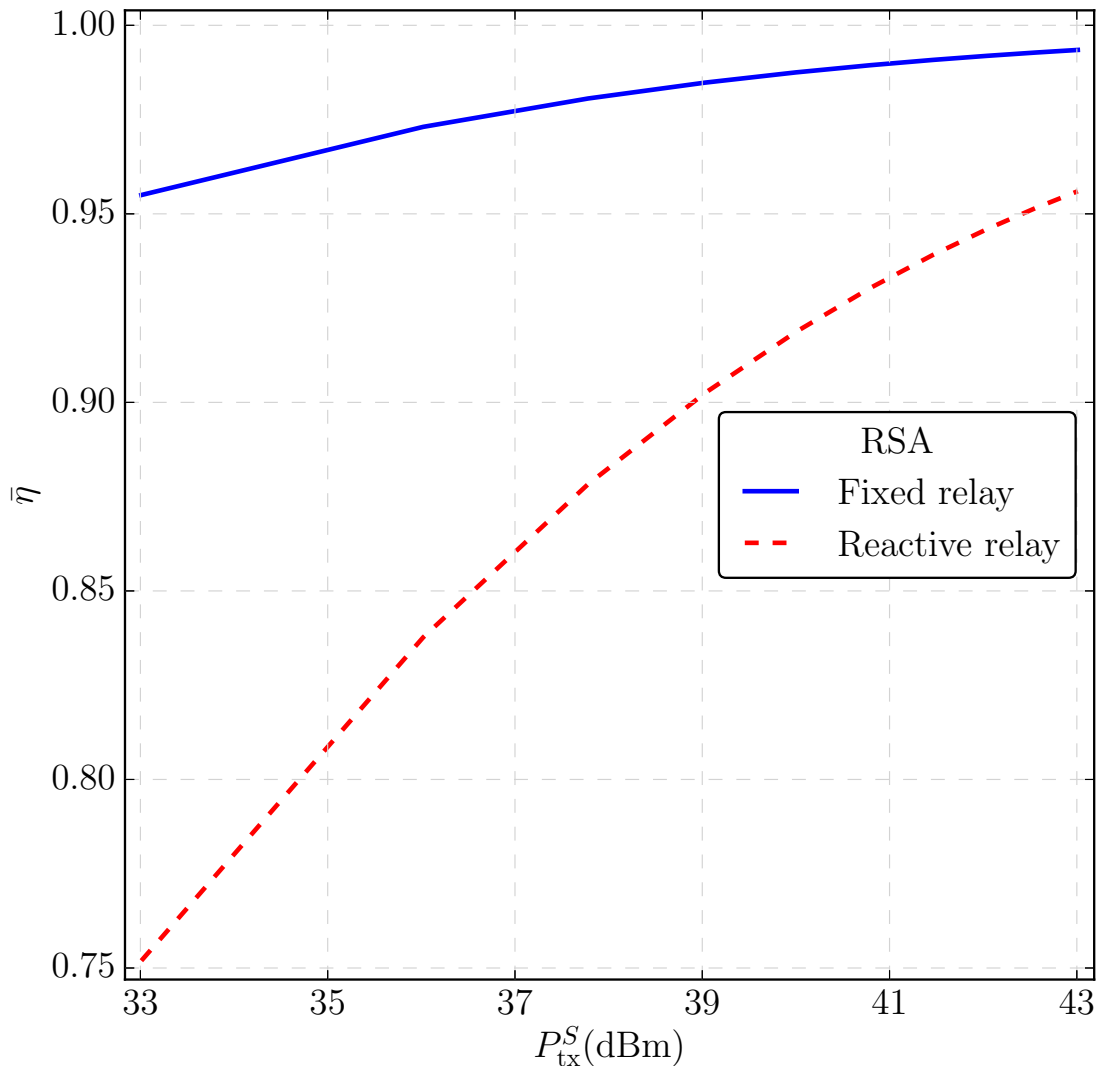


Fig. 5. Steady-state throughput efficiency increasing source's transmission power in HD configuration.

whereas interferers transmit at 30 dBm. The separation distance between source and destination is kept at 50 m and the relay is randomly located in the source's forwarding region. As expected the steady-state throughput improves with higher source's transmission power. The reactive relay benefits most from high transmission power mainly because source can reach the destination more often without undergoing long relay selection intervals to select a suitable relay.

The steady-state throughput efficiency for decreasing SI attenuation with FD configuration is shown in Fig. 6. The separation distance between source–destination is set to 50 m. As previously

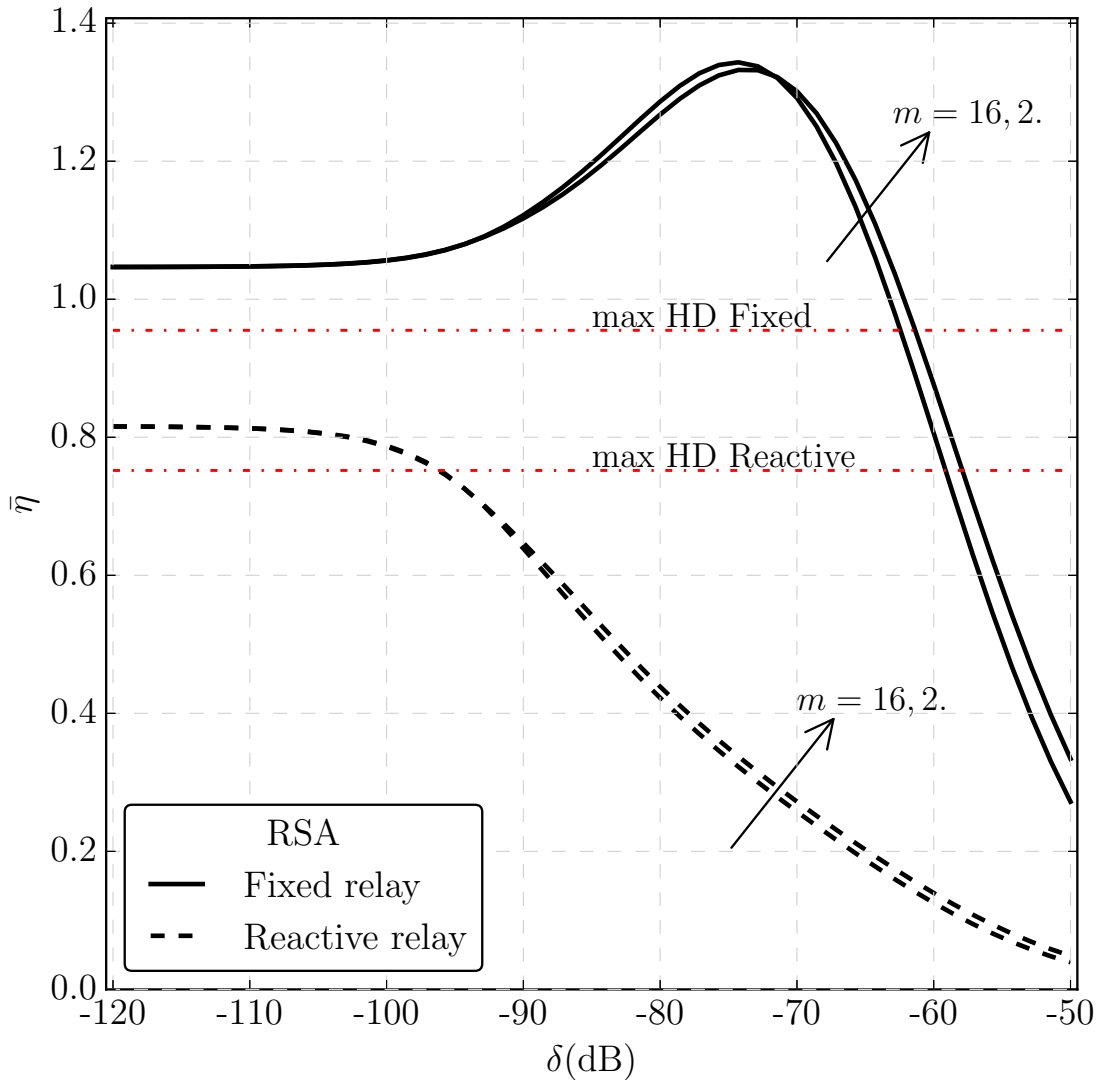


Fig. 6. Steady-state throughput efficiency for decreasing SI attenuation in FD configuration.

observed in the HD configuration, the steady-state throughput with reactive relay is severely degraded by the relay selection procedure and barely outperforms the HD scenario with 100 dB SI attenuation. It is also shown how the steady-state throughput efficiency scales for $m = 2$ and 16. As can be seen, $\bar{\eta}$ slightly degrades when the Nakagami- m parameter of the self-interference component increases from 2 (non line of sight) 16 (line of sight).

Regarding the fixed relay scheme, the FD configuration with SI attenuation ranging from 100 to nearly 70 dB show much better performance than HD mode, whereas the performance degrades

faster with lower attenuation values. From 120 to 100 dB the FD gain is negligible, because SI is low and the source node reaches the destination directly. However, for lower SI attenuation the source node uses the relay more often which increases $\bar{\eta}$ (the reward doubles). Despite that, if the attenuation is lower than 70 dB, the SI also compromises the relay performance and degrades the Steady-state throughput efficiency considerably. Moreover, as indicated by the dot-dashed threshold lines, the FD configuration outperforms the HD one for high SI attenuation values. However, if the SI attenuation drops below 70 dB, the FD performance degrades and HD becomes a better alternative. In fact, it would be advantageous to consider a dynamic forwarding scheme by which potential relays would change between FD and HD depending on the attainable SI value.

VII. CONCLUSIONS AND FINAL REMARKS

In this work, we assess how in-band full duplex relaying performs for distinct selection strategies, namely fixed and reactive. Our investigations are carried out using an analytical framework based on stochastic geometry and semi-Markov process. The former is used to model the network deployment area and radio channel, while the latter characterizes the cost of the relay selection procedure under study. Our results show that the benefits of the FD configuration are compromised by the frequent selection of the next hop relay. Moreover, the steady-state throughput degrades depending on the relative distance between source, relay and destination, as well as the respective SI attenuation. As pointed out in [14], [15], a binary search tree implementation does not take full advantage of the side information (local network topology) to shorten the relay selection time. Hence, we plan to extend this work by considering Dutch auction-based relay selection to exploit the local topology of the network in order to shorten the selection transactions. We also expect to assess the secrecy of the network, as in [33], by employing a similar framework.

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